

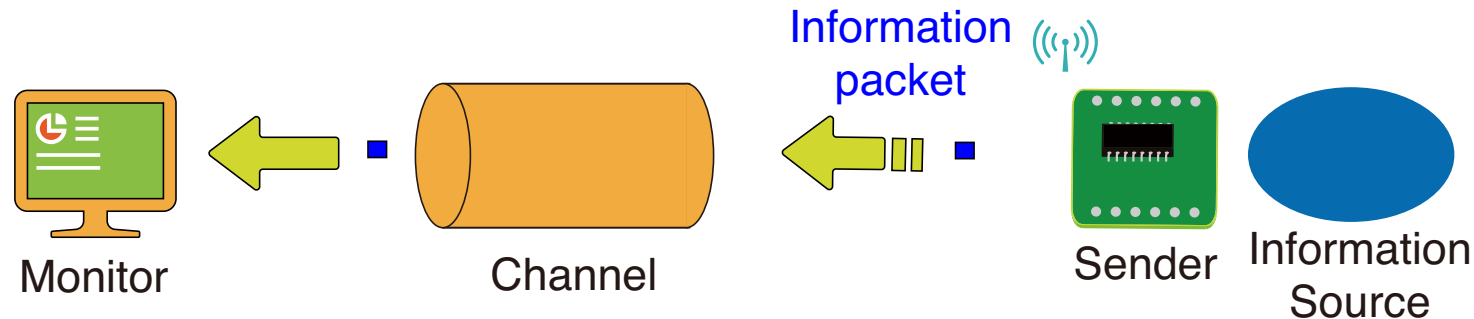
Analysis of the Stationary Distribution of the Age of Information

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Joint work with Hiroyuki Masuyama (Kyoto University)
 Tetsuya Takine (Osaka University)
 Toshiyuki Tanaka (Kyoto University)

Information Update Systems

A time-varying information source is remotely monitored



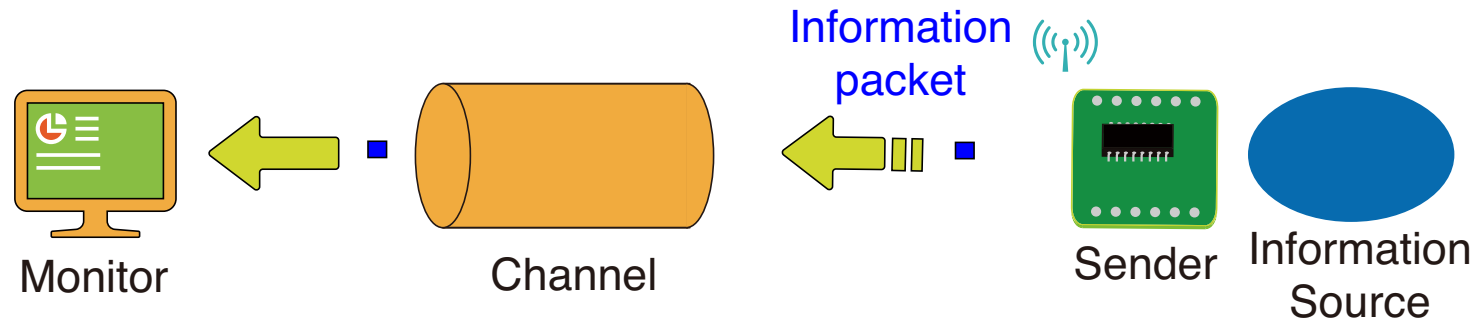
- Sender node
 - sends observed sample to the monitor
- Receiver node (monitor)
 - displays the latest information received

The displayed information is always “older” than the current state

➡ Only partial information can be obtained from the monitor

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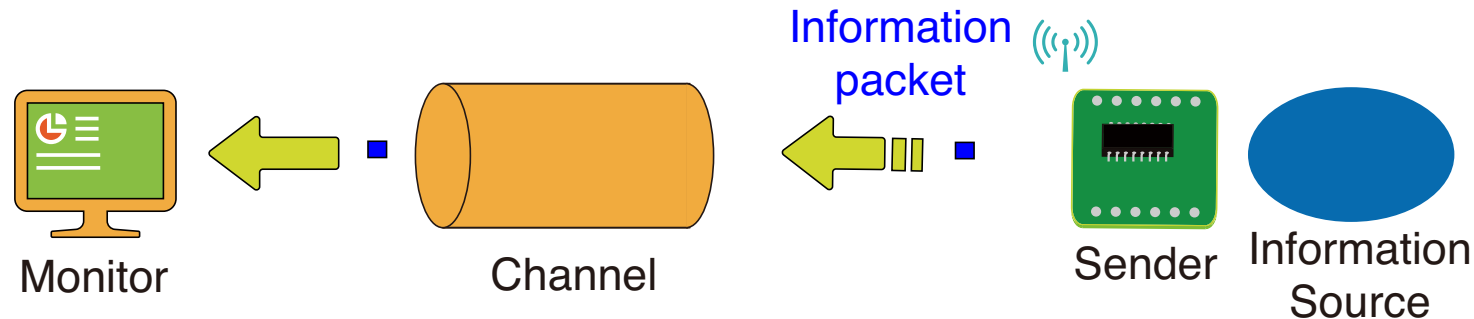
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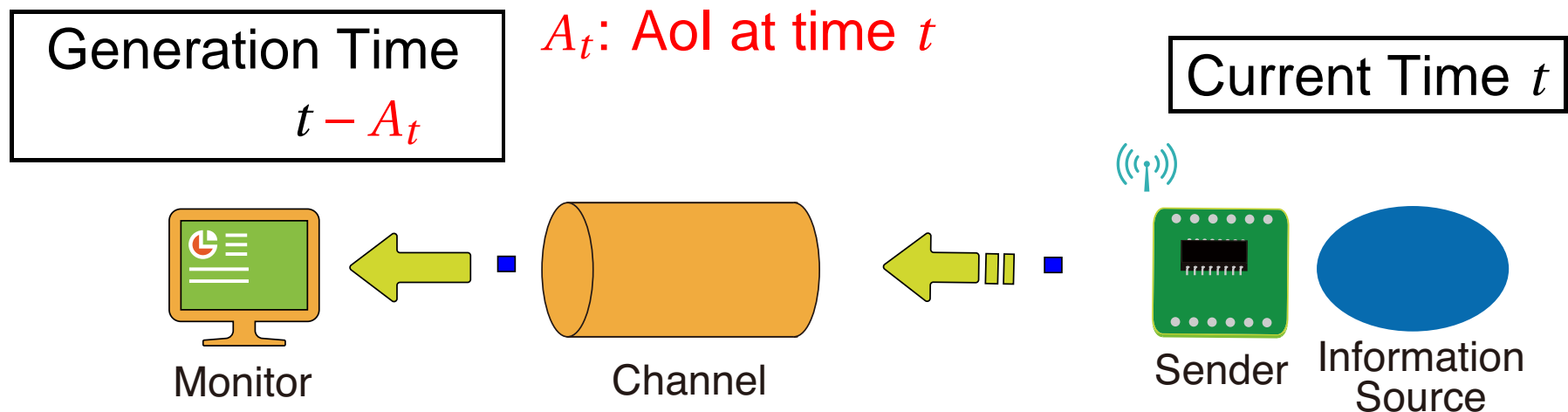
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- **Age of Information (AoI)**

A metric to characterize **the freshness** of information

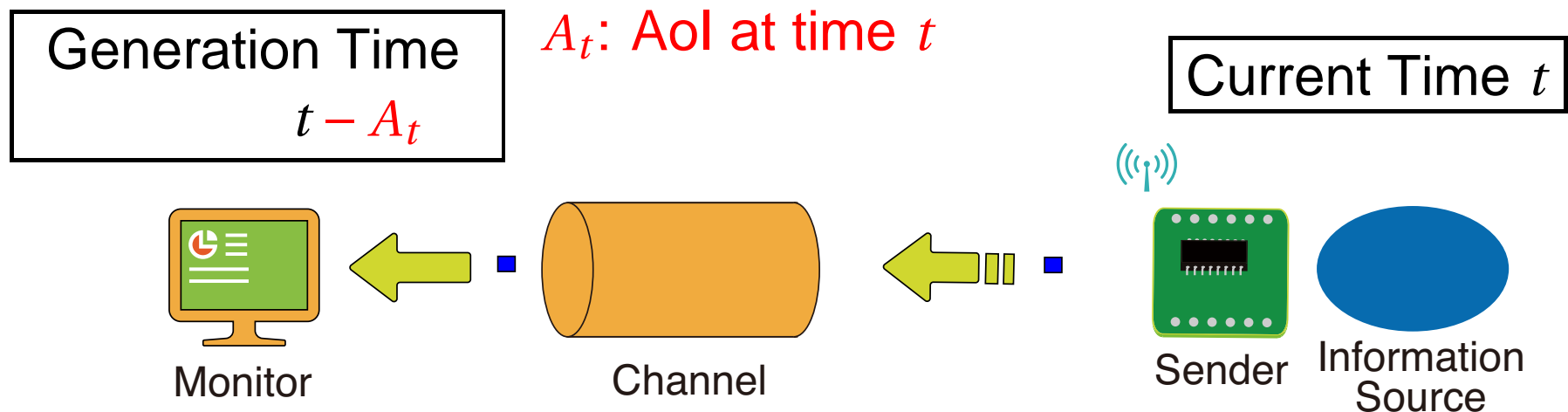
Age of Information (AoI)

- **AoI: Elapsed time of the information since its generation**
 - ◆ The smaller the AoI, the fresher the information
- No assumptions on the information source are imposed
 - ➔ The AoI defines “the freshness” in a wide class of systems



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Researches on Aol

Researches on the Aol have been expanding these years

- # of papers in IEEE international conferences
 - ◆ 2012: ISIT (1), INFOCOM (1)
 - ◆ 2013: ISIT (1)
 - ◆ 2014: ISIT (2)
 - ◆ 2015: ISIT (3), ICC (2)
 - ◆ 2016: ISIT (4), INFOCOM (1), ICC (1)
 - ◆ 2017: ISIT (12), INFOCOM (1), ICC (1), GLOBECOM (5)
 - ◆ 2018: ISIT (13), INFOCOM (2), ICC (2), GLOBECOM (4)

Outline of This Talk

- In the first half, pioneering papers on the Aol are introduced
 - ◆ Performance evaluation of VANETs [1]
 - ◆ Theoretical analysis based on the queueing theory [2]
- In the second half, our recent work is presented
 - ◆ Most previous works discuss only the mean Aol
 - ◆ We characterize the probability distribution of the Aol
 - A general formula for the Aol distribution is derived
 - An application to single-server queues is presented

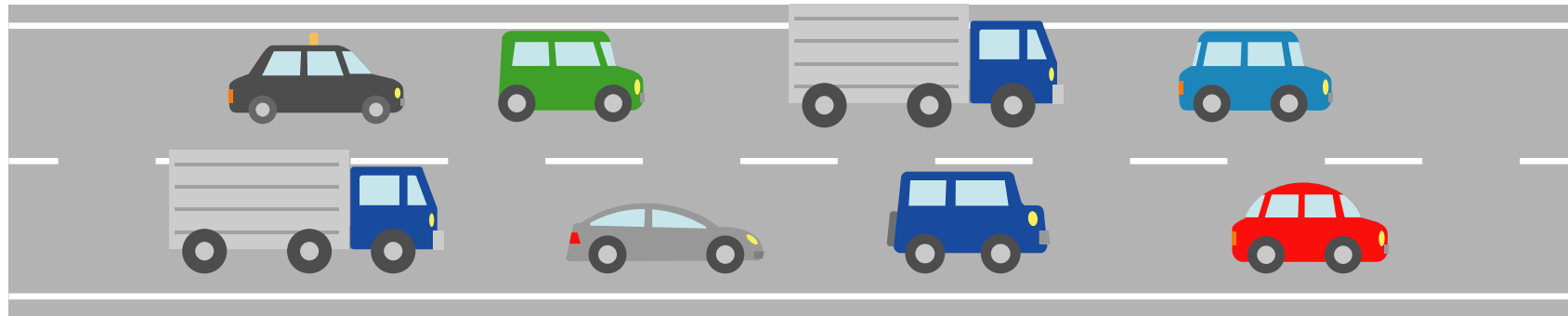
[1] S. Kaul et al., IEEE SECON 2011.

[2] S. Kaul et al., IEEE INFOCOM 2012.

Aol and Queueing Models

Vehicular Adhoc Networks (VANETs)

- Automobiles are interconnected through wireless channels
- They share each other's information to enhance driving safety
 - ◆ Its own position and driving speed
 - ◆ Information observed by a sensor and camera
 - Position and speed of neighboring cars
 - Road surface condition

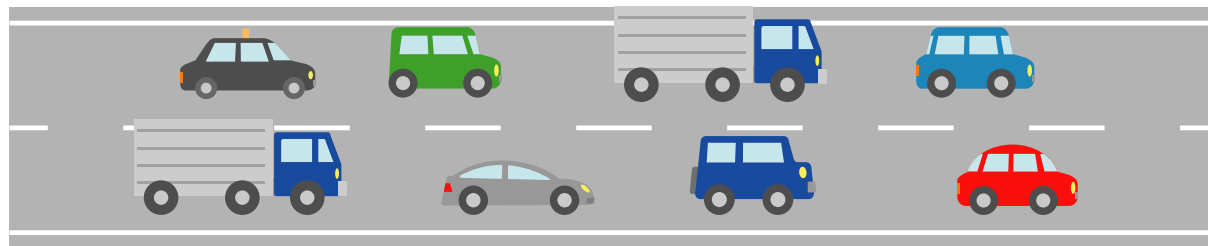


Performance measure in VANETs

- The situation may change every second
 - ◆ The value of information rapidly degenerates
 - ➔ Throughput is not an appropriate performance measure
- Freshness (Aol) is proposed as a performance measure [1]

The i th car's Aol on the j th car = $t - T_{i,j}$ (t : Current time)

- $T_{i,j}$: Time stamp of the last information i received from j

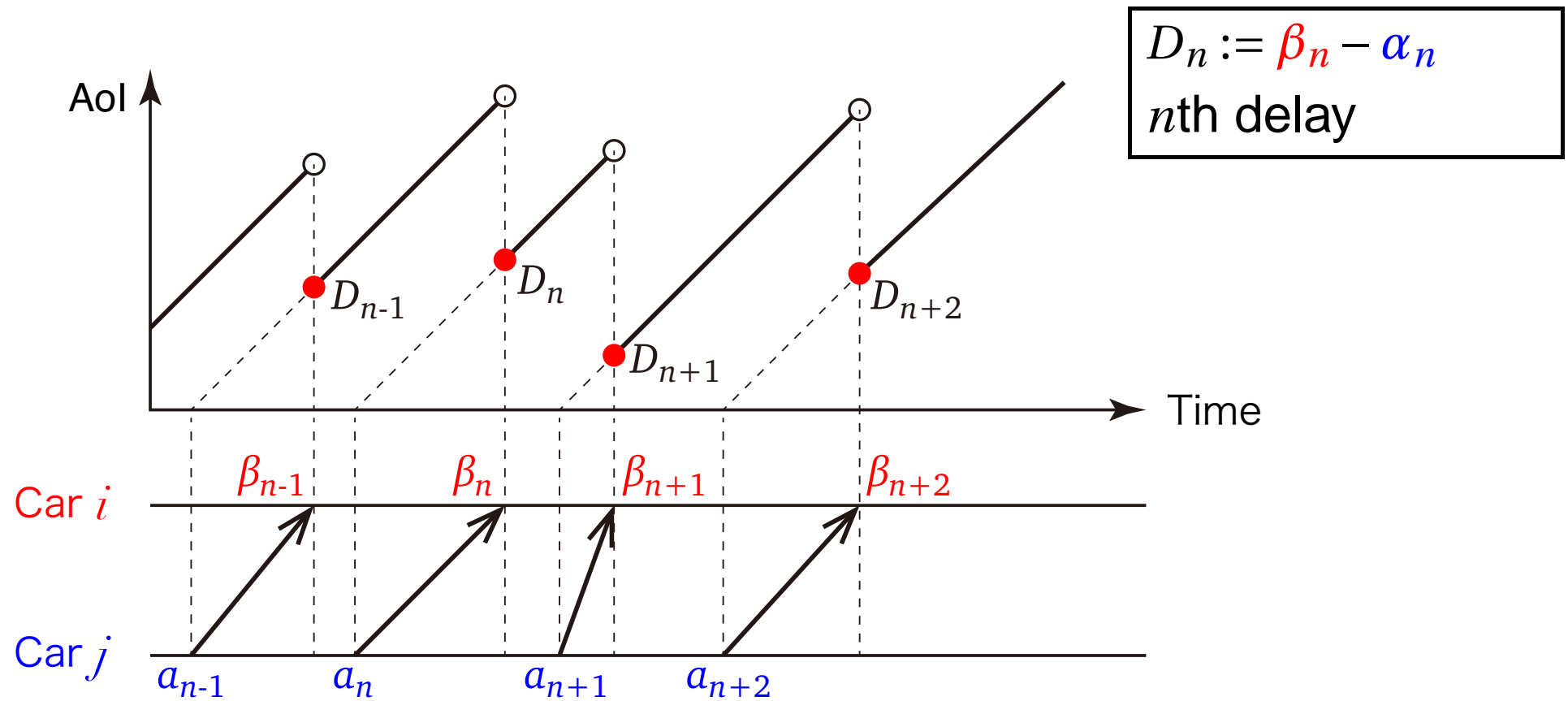


[1] S. Kaul et al., IEEE SECON 2011.

Sample Path of the Aol

Consider a specific pair (i, j) of cars

- α_n : Generation time of the n th update
- β_n : Received time of the n th update



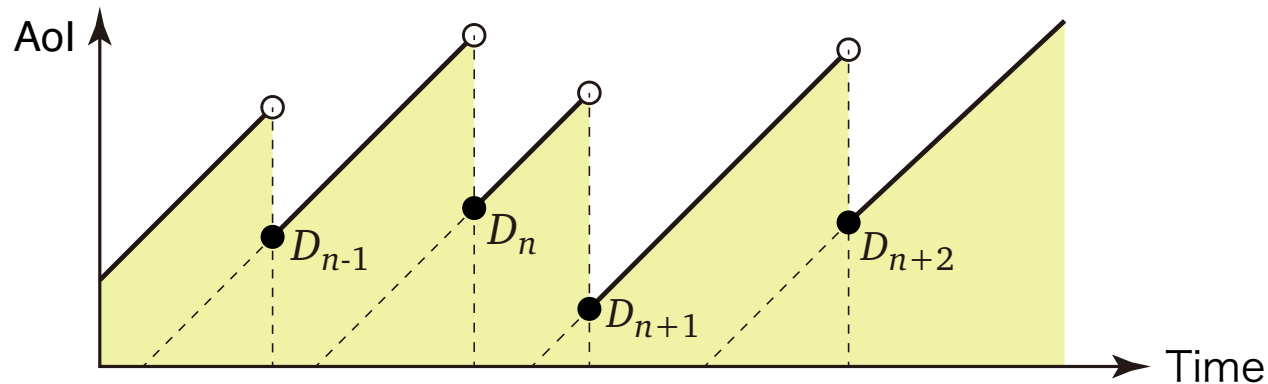
Average Aol

- In [1], the time-average of the Aol (**mean Aol**)

$$\frac{1}{T} \int_0^T A_t dt \quad \text{is proposed as a performance measure}$$

A_t : Aol at time t

- ◆ The effect of packet management on the mean Aol is evaluated with simulation experiments



[1] S. Kaul et al., IEEE SECON 2011.

Basic Properties of Aol (1)

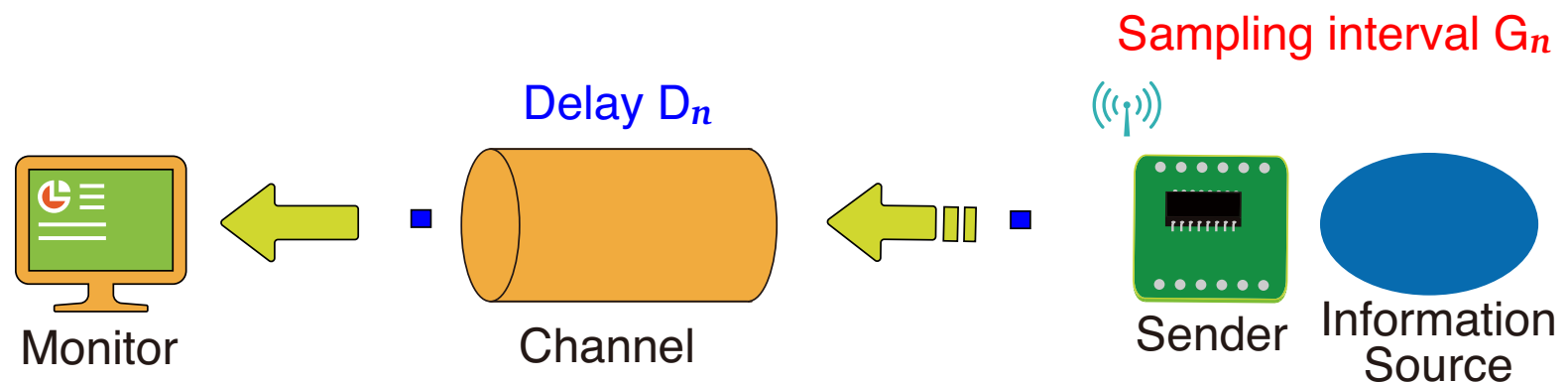
Basically, the value of the Aol is determined by

1. **Sampling interval G_n**

- ◆ If intervals are too large, information updates rarely occur

2. **Delay at communication channel D_n**

- ◆ If each packet incurs a large delay, the information cannot be kept fresh

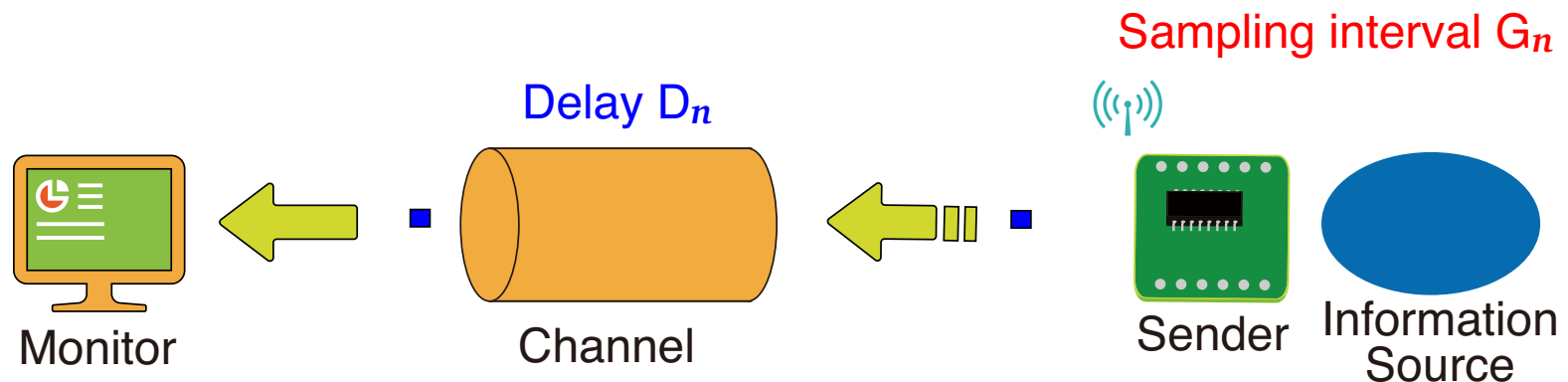


Basic properties of Aol (2)

- In order to make the Aol small,
 1. **Sampling interval G_n should be set small**
 2. **Communication delay D_n should be small**
- There is a trade-off between G_n and D_n

G_n decreases ➡ the traffic load increases ➡ D_n increases

There exists an optimal sampling interval G_n

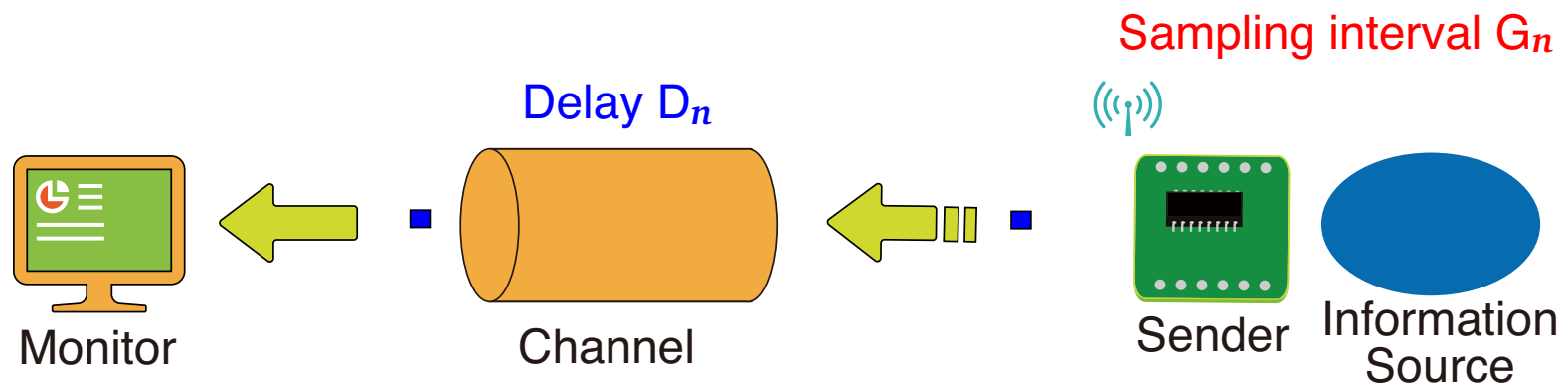


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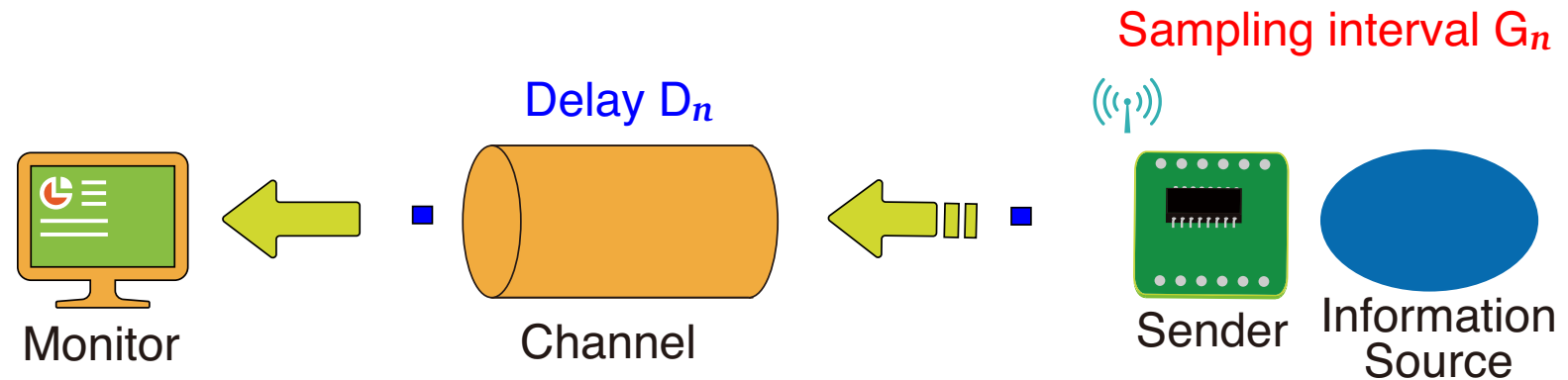
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Queueing Model and Aol

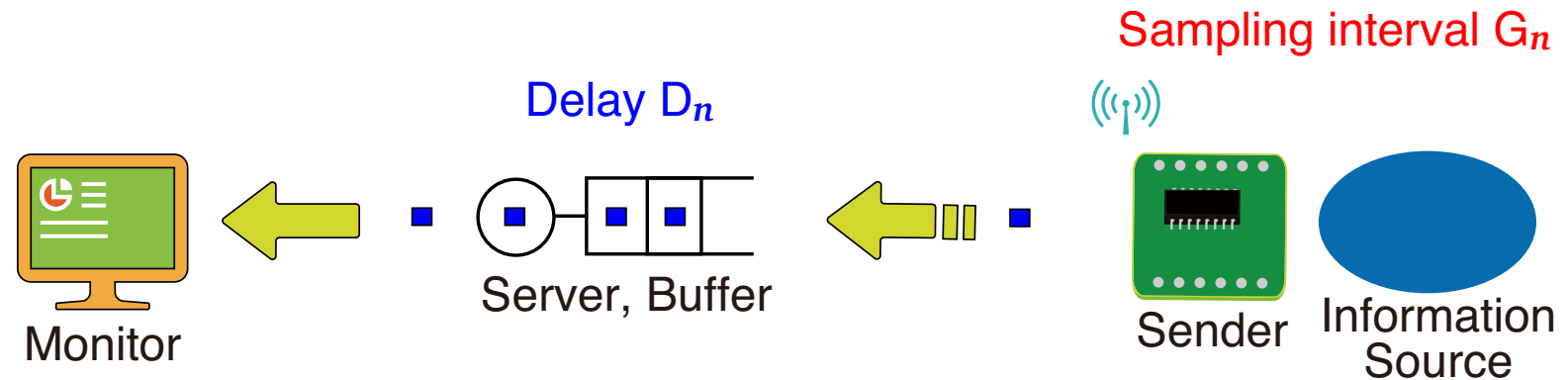
- Formulation of **the delay** using a queueing model [2]



[2] S. Kaul et al., IEEE INFOCOM 2012.

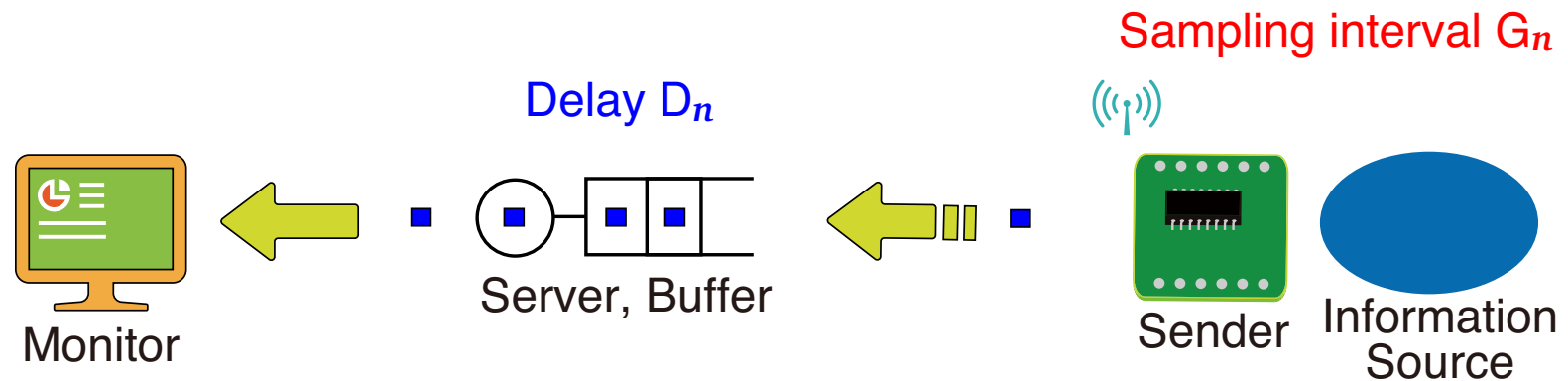
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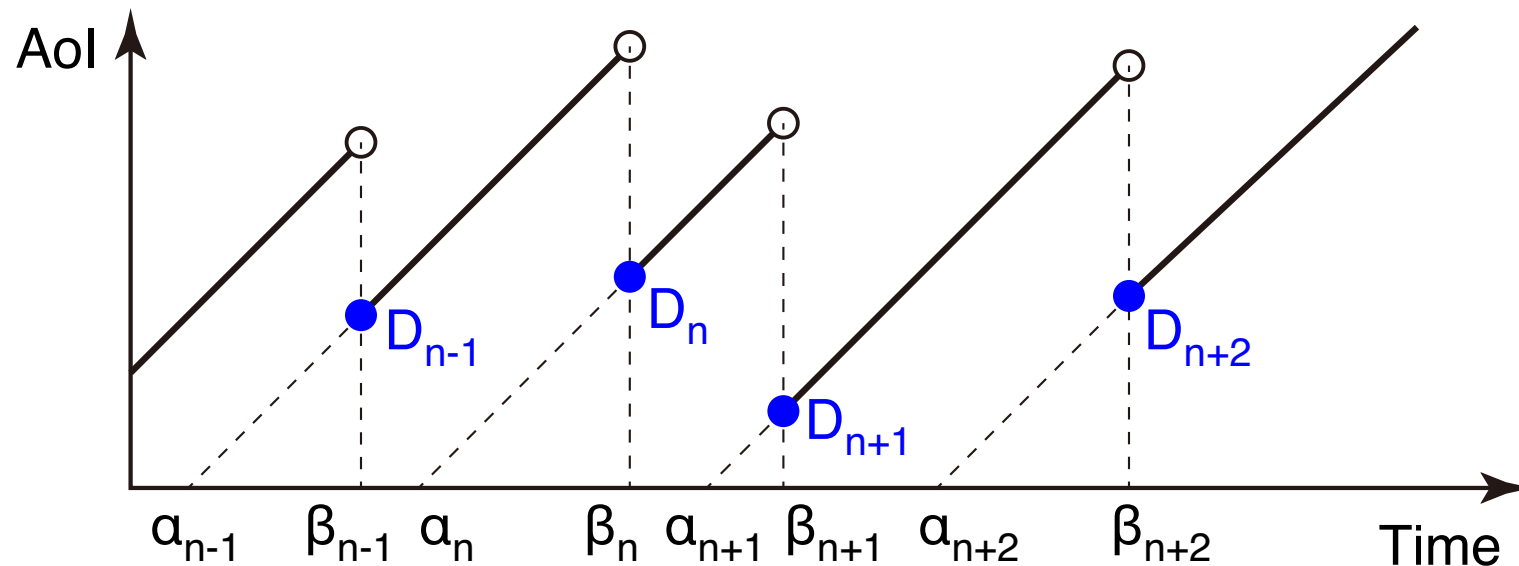
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Formula of Mean AoI (1)



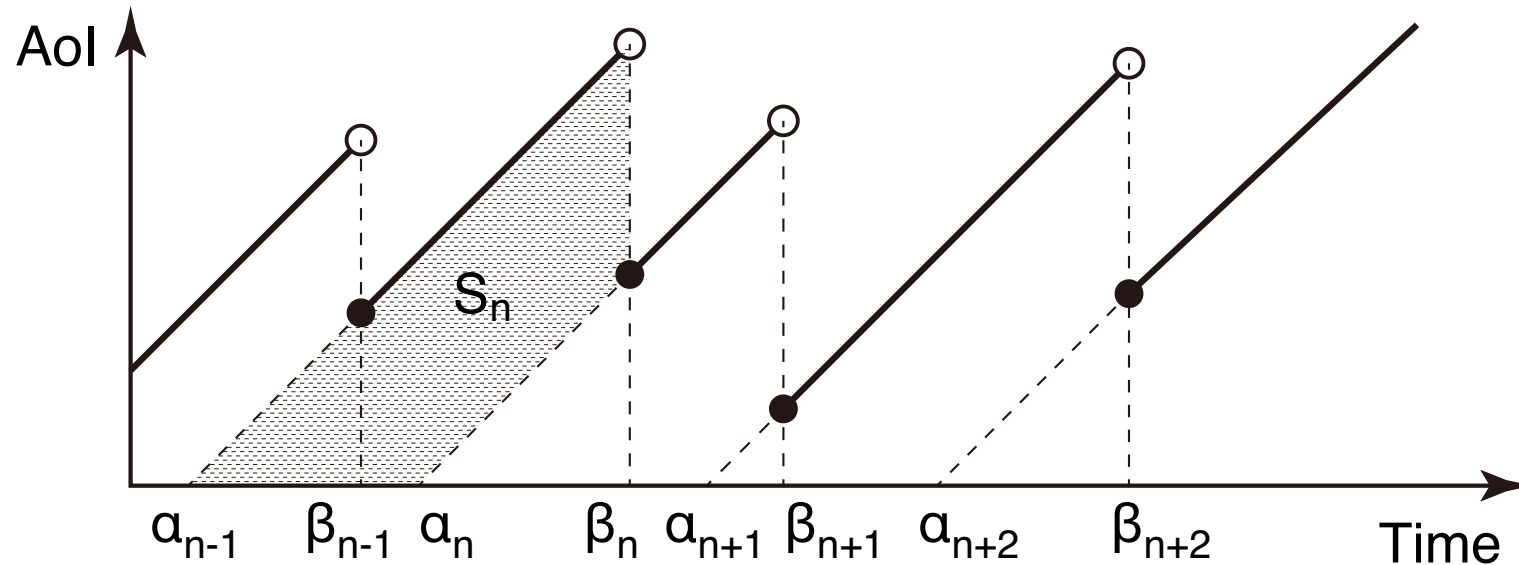
- α_n : Generation time of the n th update
- β_n : Received time of the n th update

$$G_n = \alpha_n - \alpha_{n-1}$$
$$D_n = \beta_n - \alpha_n$$



Formula of Mean AoI (2)

- The mean AoI $E[A]$ is obtained as follows [2]



$$E[A] = \lim_{T \rightarrow \infty} \frac{1}{T} \sum_{n=0}^{M_T} S_n, \quad S_n = \frac{(\beta_n - \alpha_{n-1})^2}{2} - \frac{(\beta_n - \alpha_n)^2}{2}$$

M_t : Number of received updates by time t

[2] S. Kaul et al., IEEE INFOCOM 2012.

Formula of Mean AoI (3)

$$E[A] = \lim_{T \rightarrow \infty} \frac{1}{T} \sum_{n=0}^{M_T} S_n, \quad S_n = \frac{(\beta_n - \alpha_{n-1})^2}{2} - \frac{(\beta_n - \alpha_n)^2}{2}$$

- S_n is rewritten as follows:

$$\begin{aligned} S_n &= \frac{[(\beta_n - \alpha_n) + (\alpha_n - \alpha_{n-1})]^2}{2} - \frac{(\beta_n - \alpha_n)^2}{2} \\ &= \frac{(D_n + G_n)^2 - D_n^2}{2} \\ &= \frac{G_n^2}{2} + G_n D_n \end{aligned}$$

- ◆ G_n : Inter-sampling time between $(n-1)$ st and n th updates
- ◆ D_n : Delay of the n th sample

Formula of Mean AoI (4)

G_n : Inter-sampling time between $(n - 1)$ st and n th updates

D_n : Delay of the n th sample

$$E[A] = \lim_{T \rightarrow \infty} \frac{M_T}{T} \cdot \frac{1}{M_T} \sum_{n=0}^{M_T} S_n, \quad S_n = \frac{G_n^2}{2} + G_n D_n$$

- If the system is stationary and ergodic, we have [2]

$$E[A] = \frac{\frac{E[G^2]}{2} + E[G_n D_n]}{E[G]}$$

- In general, G_n and D_n are not independent

➔ Analysis of $E[A]$ is reduced to derivation of $E[G_n D_n]$

[2] S. Kaul et al., IEEE INFOCOM 2012.

Formula of Mean AoI (4)

G_n : Inter-sampling time between $(n - 1)$ st and n th updates

D_n : Delay of the n th sample

$$E[A] = \lim_{T \rightarrow \infty} \frac{M_T}{T} \cdot \frac{1}{M_T} \sum_{n=0}^{M_T} S_n, \quad S_n = \frac{G_n^2}{2} + G_n D_n$$

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[2] S. Kaul et al., IEEE INFOCOM 2012.

Formula of Mean AoI (5)

In [2], $E[A]$ are analyzed for three queueing models

$E[H]$: Mean service time, ρ : Traffic intensity ($= E[H]/E[G]$)

$$(M/M/1) \quad E[A] = \left(1 + \frac{1}{\rho} + \frac{\rho^2}{1-\rho} \right) E[H]$$

$$(M/D/1) \quad E[A] = \left(\frac{1}{2} + \frac{1}{2(1-\rho)} + \frac{1-\rho}{\rho e^{-\rho}} \right) E[H]$$

*This explicit formula is derived in [Y. Inoue et al., IEEE ISIT 2017]

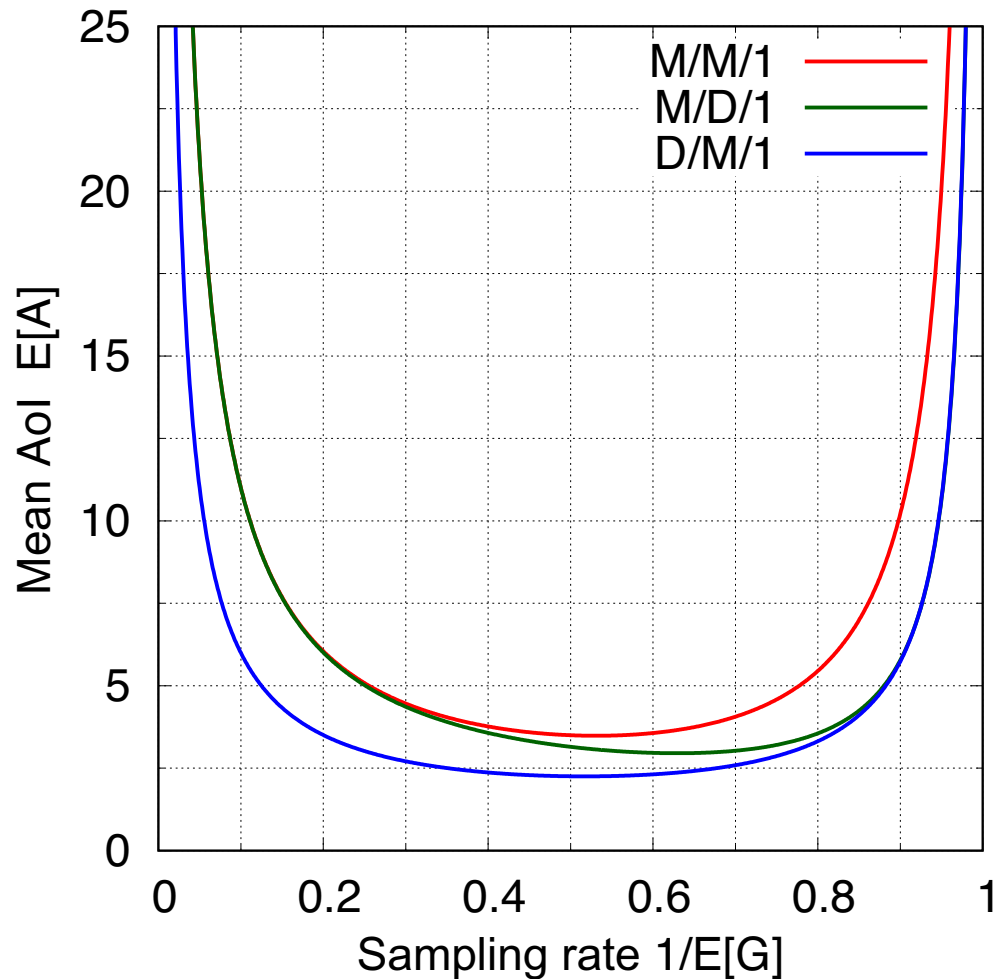
$$(D/M/1) \quad E[A] = \left(\frac{1}{2\rho} + \frac{1}{1-\gamma} \right) E[H]$$

γ is the unique solution of $x = e^{-(1-x)/\rho}$ ($0 < x < 1$)

[2] S. Kaul et al., IEEE INFOCOM 2012.

Formula of Mean AoI (6)

- We set the mean service time $E[H] = 1$



- $E[A]$ is a U-shaped function of the sampling rate

- ◆ Trade-off between the sampling rate and delay

Constant sampling interval
is efficient in terms of $E[A]$

Outline of This Talk

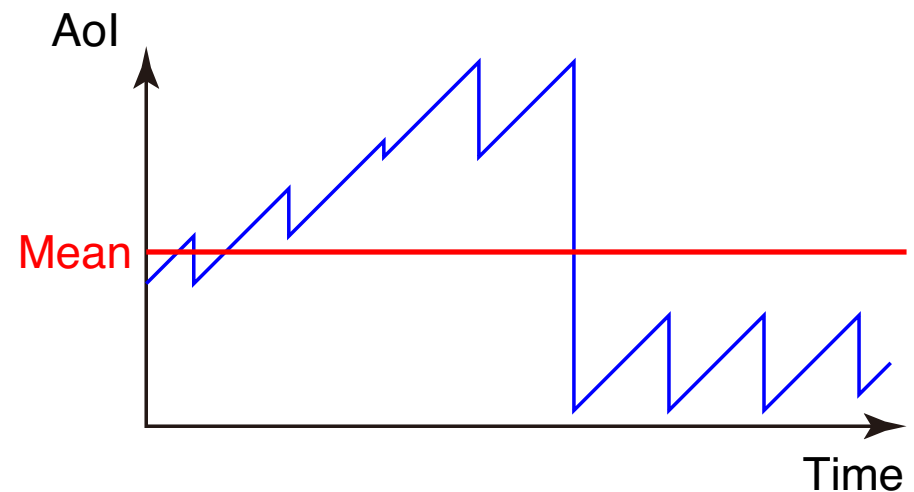
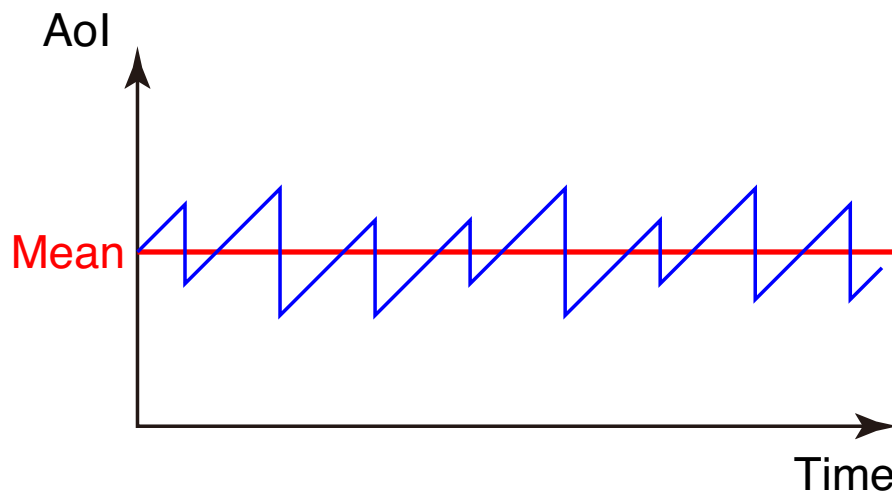
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[1] S. Kaul et al., IEEE SECON 2011.

[2] S. Kaul et al., IEEE INFOCOM 2012.

Motivation

- Although $E[A]$ is a primary performance measure, it is not sufficient to characterize the Aol process
 - ◆ In particular, **the deviation** from $E[A]$ cannot be evaluated



- We are thus interested in **the probability distribution** of the Aol

$A(x)$: Long-run fraction of time that the $Aol \leq x$

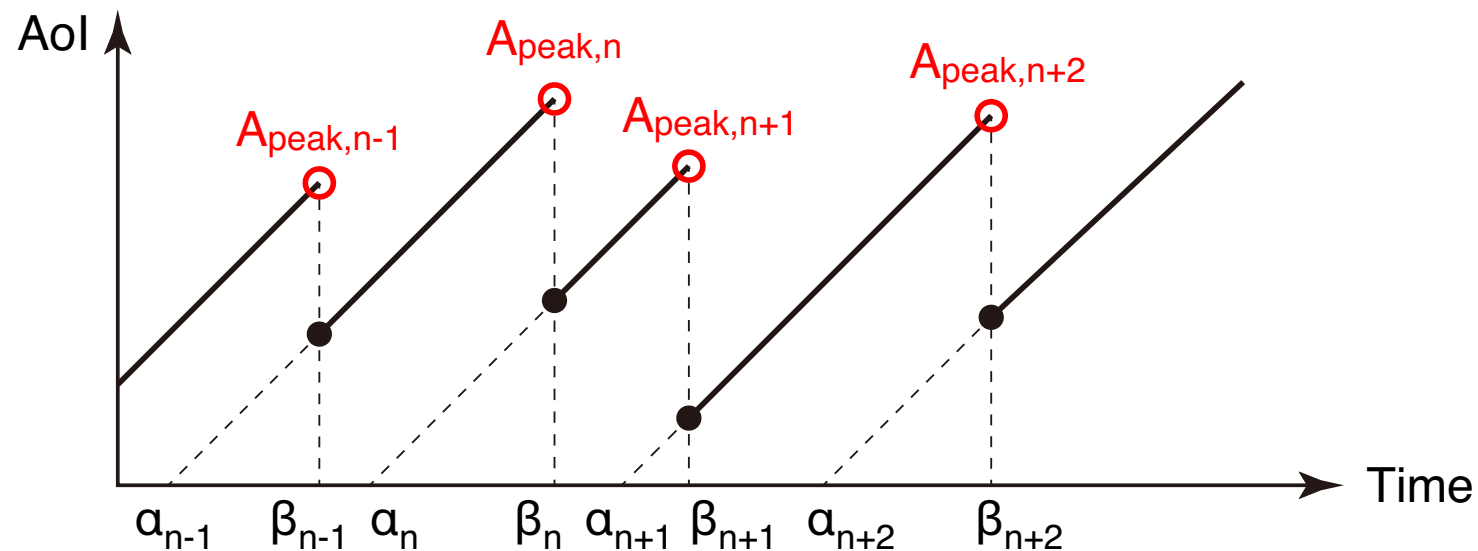
Peak AoI

- Probability distribution of the peak AoI is easier to analyze [3]
(\triangleq AoI just before an update)

$$A_{\text{peak},n+1} = D_n + (\beta_{n+1} - \beta_n)$$

D_n : Delay of the n th packet

β_n : Received time of the n th packet



[3] M. Costa et al., IEEE Trans. Inf. Theory, 62, 2016.

Outline of this work

- We derive a general formula for the Aol distribution $A(x)$:

$$A(x) = \frac{1}{E[G]} \int_0^x (D(y) - A_{\text{peak}}(y)) dy$$

$D(x)$: Delay distribution, $A_{\text{peak}}(x)$: Peak Aol distribution

- **Alternative formulas for $E[A]$** are obtained from this equation
- We present an application to the FCFS **GI/GI/1** queue
 - ◆ The distribution of the Aol is given in terms of the delay distribution
- ➔ We specialize this result to the **M/GI/1** and **GI/M/1** queues

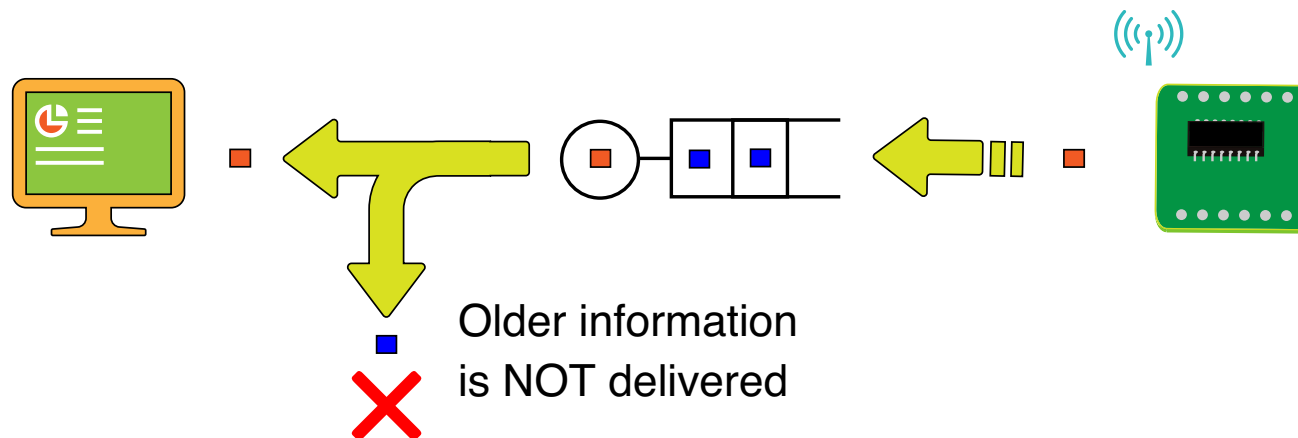
General Formula for the Aol Distribution

Classification of Information Packets

- In general, information packets are classified into two types
 - ◆ **Informative packets**, which contain **newer** information
 - ◆ **Non-informative packets**, which contain **older** information

E.g.) FCFS system ➔ All packets are **informative**

LCFS system ➔ **Non-informative** packets exist



- If we observe only the stream of **informative packets**, we have a **FIFO (First-In-First-Out)** queueing system

Sample-Path of General FIFO Queue

We thus consider a general **FIFO** queue of **informative packets**

- A **sample-path** of a general FIFO queue is characterized by

α_n ($n = 0, 1, \dots$): Arrival time of the n th packet

β_n ($n = 0, 1, \dots$): Departure time of the n th packet

α_n and β_n are **deterministic** sequences (not random variables)

- We assume the followings

(i) $\alpha_n \leq \alpha_{n+1}$ (Packets are numbered in order of arrival)

(ii) $\alpha_n \leq \beta_n$ (A packet cannot depart before its arrival)

(iii) $\beta_n \leq \beta_{n+1}$ (Packets depart in a **FIFO** manner)

(iv) $\alpha_0 \leq \beta_0 = 0 < \alpha_1$ (The system becomes empty at time 0)

Aol and Peak Aol

- M_t : Index of the last departed packet

$$M_t = \sum_{n=1}^{\infty} \mathbb{1} \{ \beta_n \leq t \}$$

$\mathbb{1} \{ X \} \triangleq \begin{cases} 1, & X \text{ is true} \\ 0, & X \text{ is false} \end{cases}$

- A_t : Aol at time t

$$A_t = t - \alpha_{M_t} \quad (\text{Current Time} - \text{Time-Stamp})$$

- $A_{\text{peak},n}$: n th peak Aol

$$A_{\text{peak},n} = \lim_{\Delta t \rightarrow 0^+} A_{\beta_n - \Delta t} \quad (\text{just before the } n\text{th departure})$$

Asymptotic Frequency Distributions

A_t : Aol at time t $A_{\text{peak},n}$: n th peak Aol

D_n : Delay of the n th packet

- $A^\#(x) \triangleq \lim_{T \rightarrow \infty} \frac{1}{T} \int_0^T \mathbb{1}\{A_t \leq x\} dt$ (The fraction of time with $A_t \leq x$)
- $A_{\text{peak}}^\#(x) \triangleq \lim_{N \rightarrow \infty} \frac{1}{N} \sum_{n=1}^N \mathbb{1}\{A_{\text{peak},n} \leq x\}$ (The relative number of peak Aols with $A_{\text{peak},n} \leq x$)
- $D^\#(x) \triangleq \lim_{N \rightarrow \infty} \frac{1}{N} \sum_{n=1}^N \mathbb{1}\{D_n \leq x\}$ (The relative number of packets with $D_n \leq x$)

We assume that these limits exist for each $x \geq 0$

General Formula for the Aol Distribution

$A^\#(x)$: Aol distribution, $A^\#_{\text{peak}}(x)$: Peak Aol distribution

$D^\#(x)$: Delay distribution

- Assumption 1: The arrival rate λ is positive and finite

$$\lambda = \lim_{T \rightarrow \infty} \frac{1}{T} \sum_{n=1}^{\infty} \mathbb{1} \{ \alpha_n \leq T \} \in (0, \infty)$$

- Assumption 2: The system is stable in the sense that

$$\lim_{T \rightarrow \infty} \frac{1}{T} \sum_{n=1}^{\infty} \mathbb{1} \{ \beta_n \leq T \} = \lim_{T \rightarrow \infty} \frac{1}{T} \sum_{n=1}^{\infty} \mathbb{1} \{ \alpha_n \leq T \} \quad (= \lambda)$$

Departure rate

Arrival rate

Under these assumptions, we have

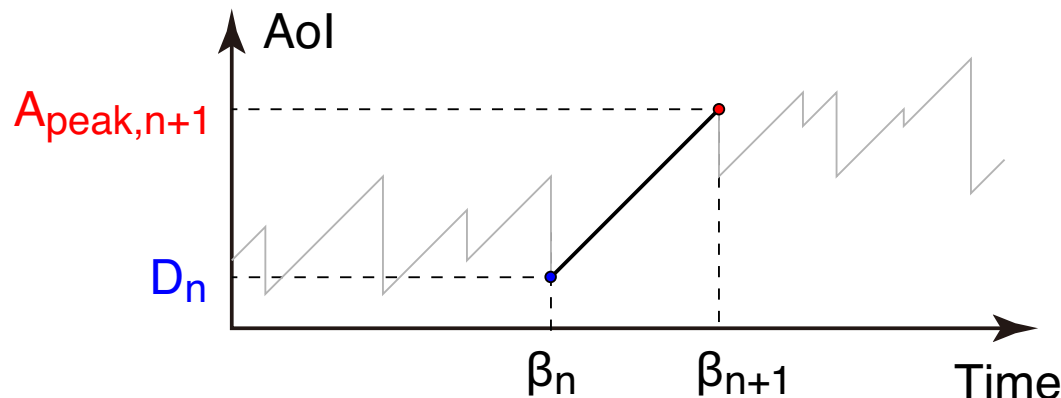
$$A^\#(x) = \lambda \int_0^x \left(D^\#(y) - A^\#_{\text{peak}}(y) \right) dy$$

Outline of Proof (1)

β_n : Departure time of the n th packet

- In a time interval $t \in [\beta_n, \beta_{n+1})$,
the Aol linearly increases from D_n to $A_{\text{peak},n+1}$
- ◆ Aol just after an information update: $A_{\beta_n} = D_n$
- ◆ Aol just before an information update: $A_{\beta_{n+1}-} = A_{\text{peak},n+1}$

$$\rightarrow \int_{\beta_n}^{\beta_{n+1}} \mathbb{1}\{A_t \leq x\} dt = \int_{D_n}^{A_{\text{peak},n+1}} \mathbb{1}\{u \leq x\} du$$

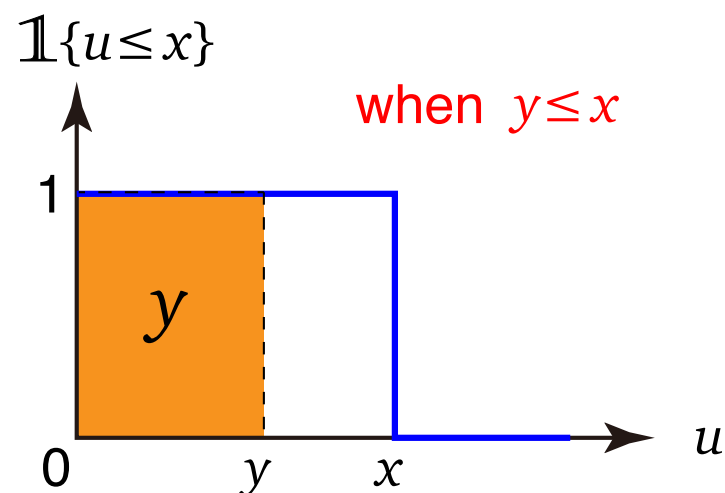
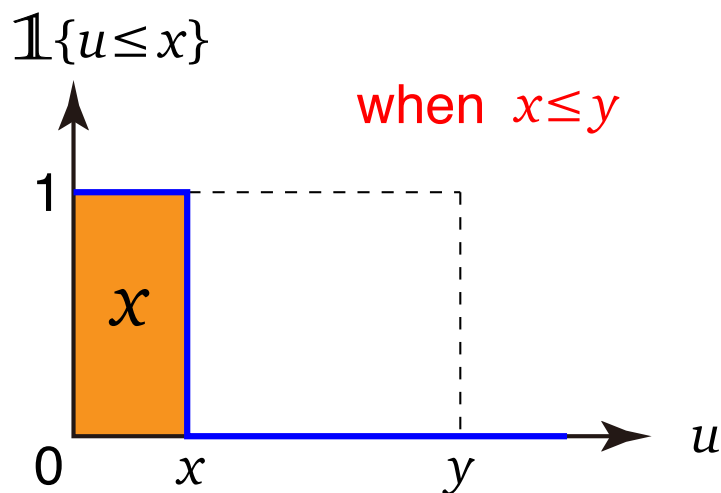


Outline of Proof (2)

$$\begin{aligned}\int_{\beta_n}^{\beta_{n+1}} \mathbb{1}\{A_t \leq x\} dt &= \int_{D_n}^{A_{\text{peak},n+1}} \mathbb{1}\{u \leq x\} du \\ &= \int_0^{A_{\text{peak},n+1}} \mathbb{1}\{u \leq x\} du - \int_0^{D_n} \mathbb{1}\{u \leq x\} du\end{aligned}$$

- For any $x \geq 0$ and $y \geq 0$, we have

$$\int_0^y \mathbb{1}\{u \leq x\} du = \min(x, y) = \int_0^x \mathbb{1}\{u \leq y\} du$$



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$$\begin{aligned}\int_0^y \mathbb{1}\{u \leq x\} du &= \min(x, y) = \int_0^x \mathbb{1}\{u \leq y\} du \\ &= \int_0^x (1 - \mathbb{1}\{y \leq u\}) du\end{aligned}$$

Outline of Proof (2)

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Therefore, we obtain

$$\int_{\beta_n}^{\beta_{n+1}} \mathbb{1}\{A_t \leq x\} dt = \int_0^x \mathbb{1}\{D_n \leq u\} du - \int_0^x \mathbb{1}\{A_{\text{peak},n+1} \leq u\} du$$

Outline of Proof (3)

$$\int_{\beta_n}^{\beta_{n+1}} \mathbb{1} \{A_t \leq x\} dt = \int_0^x \mathbb{1} \{D_n \leq u\} du - \int_0^x \mathbb{1} \{A_{\text{peak},n+1} \leq u\} du$$

- The distribution $A^\#(x)$ of the Aol is thus given by

$$\begin{aligned} A^\#(x) &= \lim_{T \rightarrow \infty} \frac{1}{T} \int_0^T \mathbb{1} \{A_t \leq x\} dt \\ &= \lim_{T \rightarrow \infty} \frac{1}{T} \sum_{n=0}^{M_T} \int_{\beta_n}^{\beta_{n+1}} \mathbb{1} \{A_t \leq x\} dt \\ &= \lim_{T \rightarrow \infty} \frac{M_T}{T} \cdot \frac{1}{M_T} \sum_{n=0}^{M_T} \left[\int_0^x \mathbb{1} \{D_n \leq u\} du - \int_0^x \mathbb{1} \{A_{\text{peak},n+1} \leq u\} du \right] \\ &= \lambda \int_0^x \left(D^\#(y) - A_{\text{peak}}^\#(y) \right) dy \end{aligned}$$

Stationarity and Ergodicity

$$A^\#(x) = \lambda \int_0^x \left(D^\#(y) - A_{\text{peak}}^\#(y) \right) dy \quad \text{holds sample-path-wise}$$

- We assume that the system is **stationary** and **ergodic**
 - ◆ **Probability distributions of system-states are time-invariant**
 - They are called stationary distributions

◆ **Stationary distributions** = **Probability distributions on a sample-path**

$$\Rightarrow A(x) = A^\#(x), \quad A_{\text{peak}}(x) = A_{\text{peak}}^\#(x), \quad \text{and} \quad D(x) = D^\#(x)$$

$A(x)$: Stationary Aol distribution

$A_{\text{peak}}(x)$: Stationary peak Aol distribution

$D(x)$: Stationary delay distribution

Notation

We use the following convention throughout the discussion below

- For any non-negative random variable F ,
 - ◆ $F(x)$: Probability distribution function of F

$$\Pr(F \leq x) = F(x)$$

- ◆ $f(x)$: Probability density function of F

$$f(x) = \frac{d}{dx}F(x), \quad x \geq 0$$

- ◆ $f^*(s)$: Laplace-Stieltjes transform (LST) of F

$$f^*(s) = \mathbb{E}\left[e^{-sF}\right] = \int_0^{\infty} e^{-sx} dF(x), \quad \operatorname{Re}(s) > 0$$

Stationary Distribution of the Aol (1)

A : Aol, A_{peak} : Peak Aol, D : Delay

- The density function of the stationary Aol

$$a(x) = \frac{D(x) - A_{\text{peak}}(x)}{E[G]}$$

- The LST of the stationary Aol

$$a^*(s) = \frac{d^*(s) - a_{\text{peak}}^*(s)}{sE[G]}$$

- The k th moment of the stationary Aol ($k = 1, 2, \dots$)

$$E[A^k] = \frac{E[(A_{\text{peak}})^{k+1}] - E[D^{k+1}]}{(k+1)E[G]}$$

Formulas for the Mean AoI

- The formula for the mean AoI $E[A]$ [1]

$$(i) \quad E[A] = \frac{\frac{E[G^2]}{2} + E[G_n D_n]}{E[G]}$$

- We have the following alternative formulas for $E[A]$:

$$(ii) \quad E[A] = \frac{E[(A_{\text{peak}})^2] - E[D^2]}{2E[G]}$$

$$(iii) \quad E[A] = \lim_{s \rightarrow 0^+} (-1) \cdot \frac{d}{ds} [a^*(s)], \quad a^*(s) = \frac{d^*(s) - a_{\text{peak}}^*(s)}{sE[G]}$$

[1] S. Kaul et al., in Proc. of IEEE INFOCOM 2012, 2731–2735, Mar. 2012.

Stationary Distribution of the Aol (2)

- The distribution $A(x)$ of the Aol is given by

$$A(x) = \lambda \int_0^x (D(y) - A_{\text{peak}}(y)) dy$$

- ◆ The Delay is widely analyzed in the queueing theory
 - ◆ We need an additional analysis on the peak Aol
- Below, we consider FCFS single-server queues
 - ◆ In the GI/GI/1 queue, $A_{\text{peak}}(x)$ is given in terms of $D(x)$
 - ◆ In the M/GI/1 and GI/M/1 queues, we can obtain formulas of the Aol from the known results for $D(x)$

Application to the FCFS GI/GI/1 Queue

GI/GI/1 Queue

- Inter-arrival times are assumed to be i.i.d. with
 - ◆ general probability distribution function $G(x)$
 - ◆ mean $E[G]$ ($\lambda = 1/E[G]$ follows)
- Processing times are assumed to be i.i.d. with
 - ◆ general probability distribution function $H(x)$
 - ◆ mean $E[H]$
- The traffic intensity $\rho \triangleq \lambda E[H]$
 - ◆ We assume $\rho < 1$ so that the system is stable
 - ◆ We also assume that the system is stationary and ergodic

Peak AoI Distribution

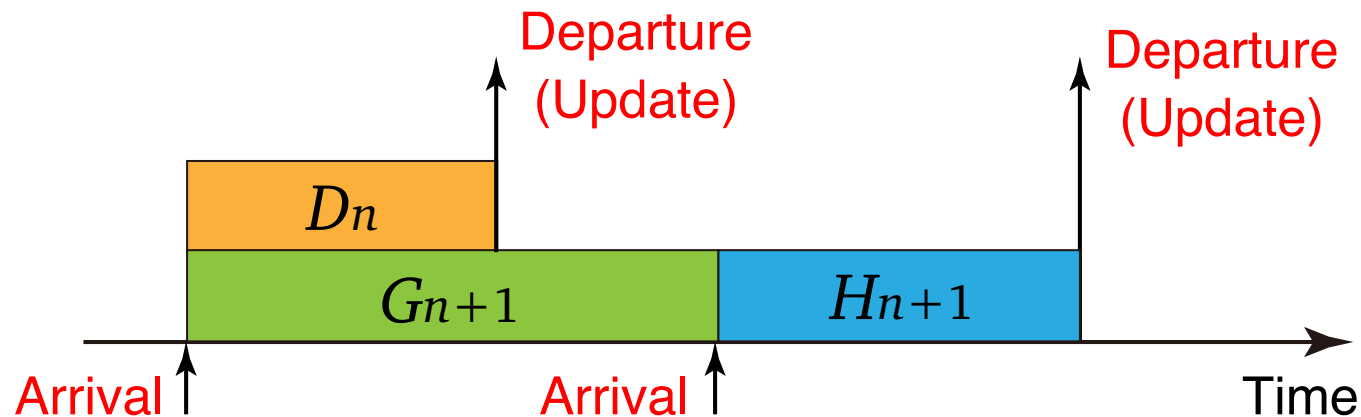
G_n : Inter-arrival time between the $(n-1)$ st and n th packets

H_n : The service time of the n th packet

D_n : The delay of the n th packet

- If $G_{n+1} > D_n$ (The system becomes empty on departure)

$$A_{\text{peak},n+1} = G_{n+1} + H_{n+1}$$



Peak AoI Distribution

G_n : Inter-arrival time between the $(n-1)$ st and n th packets

H_n : The service time of the n th packet

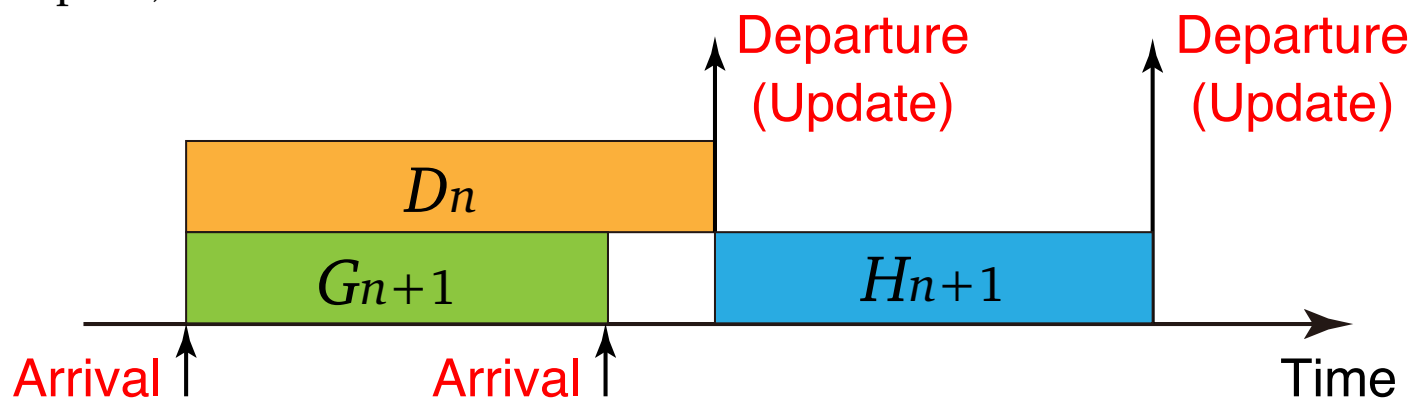
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- If $G_{n+1} > D_n$ (The system becomes empty on departure)

$$A_{\text{peak},n+1} = G_{n+1} + H_{n+1}$$

- If $G_{n+1} \leq D_n$ (The next service starts just after departure)

$$A_{\text{peak},n+1} = D_n + H_{n+1}$$



Peak AoI Distribution

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- If $G_{n+1} > D_n$ (The system becomes empty on departure)

$$A_{\text{peak},n+1} = G_{n+1} + H_{n+1}$$

- If $G_{n+1} \leq D_n$ (The next service starts just after departure)

$$A_{\text{peak},n+1} = D_n + H_{n+1}$$

We thus obtain

$$A_{\text{peak},n+1} = \max(G_{n+1}, D_n) + H_{n+1}$$

Peak AoI Distribution (2)

G_n : Inter-arrival time between the $(n-1)$ st and n th packets

H_n : Service time of the n th packet

D_n : Delay of the n th packet

- $A_{\text{peak},n} = \max(G_{n+1}, D_n) + H_{n+1}$

- G_{n+1} , D_n , and H_{n+1} are independent

➔ We obtain the LST of the stationary peak AoI

$$a_{\text{peak}}^*(s) = \left[\int_0^\infty e^{-sx} G(x) dD(x) + \int_0^\infty e^{-sx} D(x) dG(x) - \mathbb{E}[\mathbb{1}_{\{G=D\}} e^{-sG}] \right] h^*(s)$$

Summary of Results (GI/GI/1 Queue)

- The LST of the stationary Aol

$$a^*(s) = \frac{d^*(s) - a_{\text{peak}}^*(s)}{sE[G]}$$

- The LST of the stationary **peak Aol**

$$a_{\text{peak}}^*(s) = \left[\int_0^\infty e^{-sx} G(x) dD(x) + \int_0^\infty e^{-sx} D(x) dG(x) - E[\mathbb{1}_{\{G=D\}} e^{-sG}] \right] h^*(s)$$

- $a^*(s)$ is given in terms of **the delay distribution**

Special Cases: M/GI/1 and GI/M/1 Queues

M/GI/1 Queue (1)

- Exponential inter-arrival time distribution $G(x) = 1 - e^{-\lambda x}$
- General service time distribution $H(x)$ (LST $h^*(s)$)
 - ◆ M/M/1 and M/D/1 [2] are special cases of this model

In this model, the LST of the system delay D is given by

$$d^*(s) = \frac{(1 - \rho)s}{s - \lambda + \lambda h^*(s)} \cdot h^*(s)$$

➔ We obtain the LST of the stationary Aol distribution

$$a^*(s) = \rho d^*(s) \cdot \frac{1 - h^*(s)}{E[H]s} + d^*(s + \lambda) \cdot \frac{\lambda}{s + \lambda} \cdot h^*(s)$$

[2] S. Kaul et al., IEEE INFOCOM 2012.

M/GI/1 Queue (2)

- The first two moments of the Aol distribution are given by

$$\begin{aligned} \mathbf{E}[A] &= \mathbf{E}[D] + \frac{1-\rho}{\rho h^*(\lambda)} \cdot \mathbf{E}[H] \\ \mathbf{E}[A^2] &= \mathbf{E}[D^2] + \frac{2(1-\rho)}{(\rho h^*(\lambda))^2} [1 + \rho h^*(\lambda) + \lambda h^{(1)}(\lambda)] (\mathbf{E}[H])^2 \end{aligned}$$

where

$$\mathbf{E}[D] = \frac{\lambda \mathbf{E}[H^2]}{2(1-\rho)} + \mathbf{E}[H]$$

$$\mathbf{E}[D^2] = \frac{\lambda \mathbf{E}[H^3]}{3(1-\rho)} + \frac{(\lambda \mathbf{E}[H^2])^2}{2(1-\rho)^2} + \frac{\mathbf{E}[H^2]}{1-\rho}$$

GI/M/1 Queue (1)

- General inter-arrival time distribution $G(x)$ (LST $g^*(s)$)
- Exponential processing time distribution $H(x) = 1 - e^{-\mu x}$
 - ◆ D/M/1 queue [2] is a special case of this model

In this model, the system delay distribution is exponential

$$D(x) = 1 - e^{-(1-\gamma)\mu x}, \quad d^*(s) = \frac{(1-\gamma)\mu}{s + (1-\gamma)\mu}$$

where γ is a unique solution of $\gamma = g^*(\mu - \mu\gamma)$ and $\gamma \in (0, 1)$

➔ We obtain the LST of the stationary Aol distribution

$$a^*(s) = \left[\rho \cdot \frac{(1-\gamma)\mu}{s + (1-\gamma)\mu} \cdot \frac{g^*(s + (1-\gamma)\mu) - \gamma}{1-\gamma} + \frac{1 - g^*(s)}{sE[G]} \right] \frac{\mu}{s + \mu}$$

[2] S. Kaul et al., IEEE INFOCOM 2012.

GI/M/1 Queue (2)

- The first two moments of the AoI distribution are given by

$$E[A] = \frac{E[G^2]}{2E[G]} + \frac{1}{\mu} - \frac{g^{(1)}((1-\gamma)\mu)}{(1-\gamma)\mu E[G]}$$

$$E[A^2] = \frac{E[G^3]}{3E[G]} + \rho E[G^2] + \frac{2}{\mu^2} + \frac{\rho}{1-\gamma} \left[g^{(2)}((1-\gamma)\mu) - 2 \left(\frac{1}{(1-\gamma)\mu} + \frac{1}{\mu} \right) g^{(1)}((1-\gamma)\mu) \right]$$

M/M/1 Queue

- Exponential inter-arrival and service times

$$G(x) = 1 - e^{-\lambda x}, \quad H(x) = 1 - e^{-\mu x}$$

- The LST of the stationary Aol is simplified as

$$a^*(s) = \frac{(1-\rho)\mu}{s + (1-\rho)\mu} - \frac{(1-\rho)\mu s(s + \lambda + \mu)}{(s + \mu)^2(s + \lambda)}$$

- ➔ We obtain an explicit formula for the Aol distribution $A(x)$

$$A(x) = 1 - e^{-(1-\rho)\mu x} + \left(\frac{1}{1-\rho} + \rho\mu x \right) e^{-\mu x} - \frac{1}{1-\rho} \cdot e^{-\lambda x}$$

Summary

- The age of information (Aol) has been attracting a considerable attention of research communities
 - ◆ The Aol represents the freshness of information
 - ◆ Most previous works focus only on the mean Aol

In our work,

- A general formula for the Aol distribution was derived

$$A(x) = \lambda \int_0^x (D(y) - A_{\text{peak}}(y)) dy, \quad x \geq 0$$

- We presented an application to single-server FCFS queues
- A full paper including results for LCFS queues is available at <https://arxiv.org/abs/1804.06139>